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<b>APPLICATION ELEMENTS</b> <i>See MPEP chapter 600 concerning utility patent application contents</i>	<b>ADDRESS</b> <b>TO:</b> Assistant Commissioner for Patents Box Patent Application Washington, DC 20231
<p>1. <input type="checkbox"/> Fee Transmittal Form <i>(Submit an original, and a duplicate for fee processing)</i></p> <p>2. <input checked="" type="checkbox"/> Specification [Total pages 34] 31 pages specification 1 page abstract 2 pages claims 7 claims</p> <p>3. <input checked="" type="checkbox"/> Drawing(s) (35 USC 113) [Total sheets 5] <input checked="" type="checkbox"/> Informal <input type="checkbox"/> Formal [Total drawings 7]</p> <p>4. <input type="checkbox"/> Oath or Declaration [Total pages ] a. <input type="checkbox"/> Newly executed (original or copy) b. <input type="checkbox"/> Copy from a prior application (37 CFR 1.63(d)) <i>(for continuation/divisional with Box 17 completed)</i> <b>[Note Box 5 below]</b> i. <input type="checkbox"/> <u>DELETION OF INVENTOR(S)</u> Signed statement attached deleting inventor(s) named in the prior application, see 37 CFR 1.63(d)(2) and 1.33(b).</p> <p>5. <input type="checkbox"/> Incorporation by Reference <i>(usable if Box 4b is checked)</i> The entire disclosure of the prior application, from which a copy of the oath or declaration is supplied under Box 4b, is considered as being part of the disclosure of the accompanying application and is hereby incorporated by reference therein.</p>	<p>6. <input type="checkbox"/> Microfiche Computer Program (Appendix) 7. <input type="checkbox"/> Nucleotide and/or Amino Acid Sequence Submission (if applicable, all necessary) a. <input type="checkbox"/> Computer Readable Copy b. <input type="checkbox"/> Paper Copy (identical to computer copy) c. <input type="checkbox"/> Statement verifying identity of above copies</p> <p><b>ACCOMPANYING APPLICATION PARTS</b></p> <p>8. <input type="checkbox"/> Assignment Papers (cover sheet &amp; documents(s))</p> <p>9. <input type="checkbox"/> 37 CFR 3.73(b) Statement <input type="checkbox"/> Power of Attorney <i>(when there is an assignee)</i></p> <p>10. <input type="checkbox"/> English Translation of Document <i>(if applicable)</i></p> <p>11. <input checked="" type="checkbox"/> Information Disclosure <input checked="" type="checkbox"/> Copies of IDS Statement (IDS)/PTO-1449 Citations</p> <p>12. <input type="checkbox"/> Preliminary Amendment</p> <p>13. <input checked="" type="checkbox"/> Return Receipt Postcard (MPEP 503) <i>(Should be specifically itemized)</i></p> <p>14. <input type="checkbox"/> Small Entity <input type="checkbox"/> Statement filed in prior Statement(s) application, Status still proper and desired</p> <p>15. <input type="checkbox"/> Certified Copy of Priority Document(s) <i>(if foreign priority is claimed)</i></p>
16. Other: PURSUANT TO 35 U.S.C. §119, APPLICANT HEREBY CLAIMS PRIORITY TO UNITED KINGDOM PATENT APPLICATION No. 9814014.8, FILED JUNE 29, 1998	

**17. If a CONTINUING APPLICATION, check appropriate box and supply the requisite information:**

☐ Continuation      ☐ Divisional      ☐ Continuation-in-part (CIP)      of prior application No.:

☐ Cancel in this application original claims of the prior application before calculating the filing fee.

☐ Amend the specification by inserting before the first line the sentence:

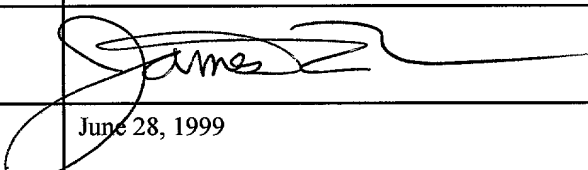
This application is a ☐ continuation ☐ divisional of application serial no. , filed , entitled , and now .

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SIGNATURE	
DATE	June 28, 1999

Title of the InventionDESIGN OF AN APPLICATION SPECIFIC PROCESSOR (ASP)Field of the Invention

The present invention relates to the design of an application specific processor (ASP) and in particular to the use of a modelling tool in the design process to aid the development of architectural modelling.

Background to the Invention

In a design process of an ASP, a model is constructed based on a CPU and a set of peripherals. Three basic types of data are required to be generated before any "real modelling" can start. These can be summarised as follows:

- A set of low level functions and constants to aid the integration of a functional model of a peripheral in a modelling language such as C.
- A set of low level functions and constants to aid the testing of the functional model. This code would execute on a simulation of the CPU and would also be useful for functional verification.
- A register or data structure mapping, indicating the size of various fields within the register, the reset state and its function.

Typically this set of data is generated by hand and must be completed before any real modelling can start. The generation by hand is a laborious task and is prone to all the usual human errors.

It is an object of the present invention to make this task less labour intensive, particularly as concerns the register or data

structure mapping.

### Summary of the Invention

According to one aspect of the present invention there is provided a method of operating a computer system to design an application specific processor (ASP) comprising:

defining a set of peripherals for the ASP which are responsive to stimuli and which communicate with a processor;

generating for each peripheral an input file which defines the functional attributes of that peripheral in a high level language with an input data structure;

entering the input file into the computer system and operating a modelling tool loaded on the computer system to generate from the input file a register definition file by allocating specific elements of the input data structure to predefined sectors of a register definition table; and

using the register definition file to create in silicon the registers of the ASP.

According to another aspect of the present invention there is provided a computer system which comprises a processor and a memory, the memory holding a program representing a modelling tool for use in designing an application specific processor (ASP), wherein the computer system comprises an input means for receiving a plurality of input files, each input file defining the functional attributes of a peripheral for the ASP in a high level language within an input data structure;

the processor being operable to execute the program representing the modelling tool to generate from the input file a register definition file by allocating specific elements of the input data structure to predefine sectors of a register definition table; and

wherein the computer system further comprises an output means for outputting the register definition file in a manner which is usable to create in silicon the registers of the ASP.

The input file can be loaded onto the computer from a physical recording device such as a disk. Likewise, the register definition file can be recorded onto a physical recording device such as a disk.

The modelling tool described herein ensures that the generation of the above sets of data are completely in sync. It also greatly reduces the time to develop the models of various peripherals by allowing the designer to concentrate on the task at hand - that of the peripheral itself. The tool also makes it possible for novice designers to get started using the peripheral modelling much quicker by not having to learn about the finer details of a particular simulator, for example by removing the need to learn details of how to send and receive data from the simulator to the peripheral.

The modelling tool is designed for use in an environment in which a simulator simulates a CPU and a memory system to which peripherals can be added.

Peripheral and subsystem support is provided by allowing the user to define a set of functions for each device that will simulate its behaviour. These functions are compiled into a device library by the user, and loaded by the simulator during execution. All the device libraries for the devices being simulated are listed in a device definition file. When the simulator is executed, these functions are dynamically linked with the simulator code.

Each device has an address range associated with it in the device definition file. Any load from or store to an address in that range will cause the simulator to execute the appropriate peripheral read or write function instead of a memory access.

The following functions are provided in each device library:

- An initialization function which is run when the simulation

starts before any application code is executed. This function must set up a structure with the names of the other functions.

- A loop function which executes regularly. The loop function is used to define asynchronous or delayed behaviour by the device, such as sending an interrupt. Each device has a loop cycle step variable which defines the frequency of execution, i.e. how many instructions are executed between two executions of the loop function. By default, the loop function is executed after every instruction.
- A function for each type of signal expected from the CPU. For peripherals this would usually be one function for Load (called the peripheral read) and one for Store (called the peripheral write). These functions are called by the simulator where an appropriate instruction is executed for an address related to the device. They define any immediate device response to the signal, such as sending back a value when a shared register location is loaded.

As an example, suppose that the application running on the CPU executes a load from a peripheral address. At this point the simulator calls the peripheral read function in the peripheral model dynamic library. The peripheral read function returns the data stored at the address and the simulator then places this value onto the stack.

A typical peripheral model and its integration within a functional simulator is shown in Figure 2. The peripheral is written in a high level language such as C and the code for this model operates upon data structures written in a manner which aids the architecture development. In order for the CPU, and hence code executing on the CPU, to access various bits of state or data, such as a control register, it must read or write to various words in memory. The read peripherals being modelled are

memory mapped in the CPU memory space. The states or data structures that are maintained within the model which are visible to the CPU, and hence to the code, must be copied to or from the registers or memory. The simulator has special calls to handle accesses to special areas such as memory.

The peripheral model writer declares an area of memory that is to be treated as special, in this case seen, as the registers memory or a data structure memory. When the CPU accesses this area of memory the peripheral model must copy or update its internal representation of the externally visible state. This is usually done by the Read and Write functions in the simulator. These functions allow the modelling to proceed in a free and easy manner, without any constraints on how it should be written or how data should be manipulated and only when it is necessary to transfer the data to the outside world is it done so via these functions. The description of the registers and data structure visible to the CPU within the peripheral will be described within the functional specification document for that peripheral.

The modelling tool described herein automates the generation of the low level functions, constants and the basis of the documentation by using the data structures specified by the peripheral model. There is sufficient information within the specification of these data structures to generate these low level functions and the documentation, using some basic conventions developed by the inventor.

These can be summarised as:

- All accesses to the data structures used within the peripheral model are done via functions. These are Query, and Set functions for each attribute (element) for each data structure.
- All accesses to the registers are done via functions with a common interface.

- The names of all the functions are derived from the attributes and structure definition of the data structures.
- All constant names used to access bits within registers are derived from the attributes and structure definition of the data structures.

For a better understanding of the present invention and to show how the same may be carried into effect reference will now be made by way of example to the accompanying drawings.

#### Brief Description of the Drawings

Figure 1 is a block diagram illustrating the modelling of an ASP;

Figure 2 is a diagram illustrating the CPU to peripheral interfaces;

Figure 3 is a diagram illustrating the function of the modelling tool;

Figure 4 is a diagram illustrating how the files are derived from the data structure of the input file using the modelling tool;

Figure 5 is a flow chart illustrating the high level operation of the modelling tool;

Figure 6 is a sketch illustrating the use of the files in a simulation; and

Figure 7 is a sketch illustrating how a functional model can interact with a real simulation.

#### Description of the Preferred Embodiments

An application specific processor is modelled as a central processor (CPU) 2 and a set of peripherals 4. The CPU 2 is modelled with the basic elements of an interrupt handler 6 and memory 8. A set of applications running on the CPU are denoted by the process circle 10 labelled APPLS. Each peripheral 4 is modelled with an internal interface 12 between the peripheral and



the CPU 2 and an external interface 14 between the peripheral and the "outside world", that is externally of the ASP. At the time of modelling the ASP, it is not known whether or not the peripherals will in fact be implemented in software, hardware or some combination of both. However, whether finally implemented in software or hardware or some combination of both, the peripherals 4 represent how the central processor 2 cooperates with the external environment. The external interfaces 14 receive stimuli S from the external environment and generate responses R in response to the stimuli. These are carried by the external interface 14. The internal interfaces 12 carry state information and data between the peripherals and the applications 10 running on the CPU 2. This is described in more detail with reference to Figure 2.

Figure 2 illustrates a single peripheral 4 which is to be modelled as a plurality of peripheral processes P1,P2,P3 etc. The CPU 2 is shown only with its applications 10 and a register bank 16. The register bank 16 represents a particular area of memory which may be modelled as registers or a conventional RAM which is dedicated to activities of that peripheral. A number of different applications may be running on the CPU 2, denoted APP1,APP2, etc. The applications 10, running on the CPU are able to write data to the register bank 16 along the write path 18 and can read data from the register bank 16 along the read path 20. These register read and writes are simulated as CPU Read/Write functions. In addition the peripheral 4 needs to be able to receive data and state from the register bank 16 and to write data and state to the register bank 16. This is accomplished by the interface 12. The modelling tool described herein is valuable for implementing the interface 12 in the modelling phase, the simulation phase and the implementation (silicon) phase of the design process. It is not a trivial matter to model, simulate or implement the interface 12. In designing an ASP, the peripherals 4 are modelled in a high level language such as C. One of the facets of that language is that the data structure which is utilised is written in a manner which aids

architecture development in particular in terms of its portability between environments. In particular, it allows the definition, modification and access of individual elements very easily, regardless of the length of the element. This is a particularly useful feature when designing or modelling because it means that the length of elements can be altered without requiring a complete revision of the modelling code. C also allows very simple access to an element, regardless of its length. However, this facet of C and other high level languages creates a practical difficulty when the code developed in that language has to be simulated with applications running on a conventional CPU and using fixed length registers. The tool described herein provides a mechanism for greatly simplifying this task.

Figure 3 shows in general terms how this is achieved.

An input file is created for each peripheral 4 in a high level language such as C using an input data structure compatible with that language. That input file defines the interfacing behaviour of the peripheral 4 with respect to the CPU. The architect determines the responses R of the peripheral with respect to external stimuli S. A modelling tool 24 generates automatically from the data structure defined in the input file 24 a documentation file 26, an interface functions file 28, and a test functions file 30.

The interface functions file 28 contains a set of "glue" functions which are derived from the individual elements of the data structure in the input file 22 but which are defined in a manner which is independent of any particular data structure. The "glue" functions define the attributes of the interface 12 and include:

- constant definitions
- read functions
- write functions

query functions  
set functions.

The constant definitions define the context for the peripheral.

In particular, they define the address range in memory associated with the device being modelled by that peripheral and bit locations within the registers of particular elements of the data structure. Any load from or store to an address in that defined range will cause a simulator to execute the appropriate peripheral read or write function instead of a memory access.

The read and write functions allow the peripheral to read and write data and state from and to the specified register of the CPU.

Query functions allow the peripheral to request a value from a specified register in the CPU.

Set functions allow the peripheral to write a value to a specified register of the CPU.

The documentation file defines the registers and their contents for use in setting up a simulator on the CPU.

The test functions take the form define the attributes of the CPU read/write paths 18, 20 and include:

constant definitions  
read functions  
write functions.

The constant definitions match those already defined as part of the interface function file 28. Likewise, the read and write functions allow the CPU to read and write from the specified registers. Once again the functions are defined such that they have a common name but are implemented in a manner which is



Read From NAME

Write To NAME

In the example of Figure 4, this is done for both the SarControlRegister and the SarSegmentationContextRegister. The query and set functions are defined by reference to each ELEMENT E of the data structure in the form:

ELEMENT In NAME

Set ELEMENT In NAME

The documentation file 26 is set up for each register by deriving information directly from the data structure as indicated in Figure 4. Thus, each register definition comprises the following parameters:

word offset	-	defining an offset location of the register in memory
bit offset	-	defining the bit location of each element in the register and derivable from the bit length BL in the data structure
bit field	-	naming the element of that bit location
function	-	defining the function F of the element
reset state	-	value of entity on reset
read/write	-	whether entity read or writable from CPU

The contents for each field to define these parameters can be derived directly from the data structure of the input file 22. To avoid over-complicating the figure, the arrows are shown only for the read and write functions in respect of the SarControlRegister and, as far as the documentation file is concerned, only for the first bit location of that register. Tables 1 and 2 show the complete documentation files for the SarControlRegister and SarSegmentationContextRegister.

For each of the typedefs in the input file a table will be

generated which will describe the allocation of the attributes to the words that make up the data structure in the CPU memory space. Each table will also describe the allocation of the bits within the word(s) as well as the meaning associated to these bits. The reset state will be given, and whether the attribute (bits are read, writable or both. The allocation of the bits within a word and indeed the words themselves will be driven by command line arguments to the modelling tool. The documentation file can be output in various formats, for example ascii and mif. The files are intended to be included or pasted into the main functional (or other) specification of the peripheral.

Figure 5 is a flow chart illustrating high level operation of the modelling tool 24. At step S1, input parameters given to the modelling tool are checked. At step S2, the input file corresponding to one of the specified parameters is opened. Its contents are checked and any areas are reported in a meaningful manner (step S3) by an error routine. If the contents of the input file are valid, the files are opened and named at step S4. At step S5, the files are created as described earlier with reference to Figures 3 and 4. Finally, at step S6, the files are closed.

Some specific examples are given in the following annexes.

Annexe 1 is an exemplary BNF sequence (Backus Naur Form of notation) for an input file 22. Annexe 2 is an example of a simple data structure within the input file, and Annexe 3 is an example of a data structure of medium complexity within the input file.

Annexe 4 is an exemplary BNF sequence for the read function of the interface functions file for a data structure of medium complexity and Annexe 5 is an example of an output fragment.

Annexe 6 is an exemplary BNF sequence for a write function for the interface functions file for a data structure of medium

complexity and Annexe 7 is an example of an output fragment.

Annexe 8 is an exemplary BNF sequence for a query function for a data structure of a simple type and Annexe 9 is an example output fragment. Annexe 10 is an exemplary BNF sequence for a set function of a simple data structure type and Annexe 11 is an example of an output fragment.

For the test functions file 30, Annexe 12 is an exemplary BNF sequence for a read function for a data structure of medium complexity, and Annexe 13 is an exemplary output fragment. Annexe 14 is an exemplary BNF sequence for a data structure of medium complexity for the write function of the test functions file 30 and Annexe 15 is an exemplary output fragment.

Annexe 16 is one example in BNF format of a documentation file.

Figure 6 shows how the modelling tool used herein is used in the simulation phase of a design process. Each peripheral device has a set of functions which simulate its behaviour. These are created as the input file 22 for each peripheral device. As already explained, each peripheral device may have a number of different simulated processes, P1, P2, P3, etc (see Figure 2). The input file 22 defines each peripheral device and therefore may include information about each of the individual simulation processes. As described above, the input file is entered into a processor which is loaded with the modelling tool 24 and which thus generates the various files illustrated in Figure 3. As already mentioned, these include an interface functions file 28 and a test functions file 30. The interface functions file for each peripheral device is held in a device library 40. In Figure 6, the denotation IFP1 denotes the interface functions for the simulated process P1 of the peripheral device. The test functions for each simulated process form part of the simulation code for the applications to be run on the CPU. The denotation TFP1 denotes the test function for the simulated process P1. In Figure 6 it is illustrated as running in the simulated version

of application 1, APP1. A device definition file 42 holds a list of the device libraries. Each device has an address range associated with it in the device definition file. Any load from or store to an address in that range will cause the simulator to execute the appropriate peripheral read or write function instead of a memory access. For example, if the simulator processor attempts to access an address in range A0 to A3, this maps onto the device library 40 holding the simulating processes for the peripheral of Figure 2. Instead of allowing the access to go ahead, that causes the simulator processor to perform the function defined by the relevant interface function file. This causes data to be written to, accessed from or modified in the representation of the register bank 16 associated with that peripheral. This type of access may have been caused by the simulator processor running code from the test functions incorporated within the process being simulated, APP1 or by code within the device libraries if the peripherals are at that time being simulated. Either way, the representation of the register bank 16 associated with that peripheral device is kept correctly updated.

The modelling tool described herein gives rise to another advantage. Figure 7 illustrates a functional model for the application specific processor (ASP) running in a high level language such as C, and a real simulation which will run in a simulation language such as VHDL. The real simulation simulates the actual implemented chip down at the transistor level, and is used before the circuit which has been designed is actually implemented in silicon. The real simulation itself is necessary prior to implementing a circuit in silicon to try and establish as far as possible how an actual device will operate. However, real simulations are very slow. Conversely, the functional model itself can run quite quickly, although of course that is only modelling the architecture and not the actual silicon implementation as designed. However, because the modelling tool described herein generates matching test function and interface function files for each peripheral device, it is possible to



speed up the real simulation by running the functional model for an initialisation or set up phase (or any other phase) and, at the end of that phase, extracting the state of the application programs, APP1, APP2 at a particular point in time and the state of the peripheral devices at the same point of time. The state of the application programs and the environmental stimuli can be derived from the test function files 30 and the state of the peripheral devices can be derived from the interface function files 28. The contents of these files are loaded into a modelling file which is translated into a simulation file which can be loaded into the real simulation process. The modelling file can be in memory or on disk, as can the simulation file. Thus, it is possible for the functional model and the real simulation to run on the same CPU, with a transfer of the modelling file to the simulation file in the CPU memory. Alternatively, the functional model can be run on one CPU, with the modelling file being on a transferable disk which can be translated onto a simulation file and loaded into another CPU running the real simulation. The extraction of state from the functional model at a particular point in time in order to "kick start" the real simulation significantly reduces the overall simulation time. The environmental stimuli from and to the peripheral devices at that point of time can also be derived from the functional model and loaded into the real simulation.

Another advantage of the modelling tool described herein is its generation of the documentation file 26. This defines the actual registers and can be used therefore to implement these registers in a final silicon implementation. This significantly reduces the amount of manual design work that needs to be carried out.

TABLE 1. SarControlRegister

Word Offset	Bit Offset	Bit Field	Function	Reset State	R/W
0	0	StartSeg-mentation	Enables seg-mentation engine	0	R/W
	1	EnablePac-ingEngine	Enables pac-ing engine	0	R/W
	2	EnablePac-ingClock	Enables pac-ing clock	0	R/W

TABLE 2. SarSegmentationContextRegister

Word Offset	Bit Offset	Bit Field	Function	Reset State	R/W
0	0:31	Value	Context start address	0	R/W

ANNEXE 1

```

InputFile ::= Typedefs
Typedefs ::= Typedef {Typedef}
Typedef ::= <TypedefToken> <StructToken> TypedefStructureName TypedefBody
TypedefBody ::= <OpenBraceToken> AttributeDefinitions <ClosingBraceToken> TypedefIdentifierName <SemiColonToken>
AttributeDefinitions ::= AttributeDefinition {AttributeDefinition}
AttributeDefinition ::= AttributeType AttributeIdentifierName
                        {<ColonToken> AttributeTypeBitSize} <SemiColonToken> {AttributeCommentsField}
AttributeType ::= ValidAttributeType
AttributeCommentsField ::= {AttributeReadWriteComment} AttributeFunctionComments
ValidAttributeType ::= BoolType | UInt32Type | St20wordType |
                        Ust20wordType | ByteType | UByteType |
                        Int16Type | UInt16Type | UWordType | WordType |
                        {OtherType}
AttributeReadWriteComment ::= <OpenCStyleCommentToken> ReadWriteSelector <CloseCStyleCommentToken>
ReadWriteSelector ::= ReadSelect | WriteSelect | ReadWriteSelect
AttributeFunctionComments ::= AttributeFunctionComment {AttributeFunctionComment}
AttributeFunctionComment ::= <OpenCStyleCommentToken> FunctionSpecifier <CloseCStyleCommentToken>

BoolType ::= BOOL
UInt32Type ::= UINT32
St20wordType ::= ST20WORD
Ust20wordType ::= UST20WORD
ByteType ::= BYTE
UByteType ::= UBYTE
Int16Type ::= INT16
UInt16Type ::= UINT16
UWordType ::= UWORD
WordType ::= WORD
CloseCStyleCommentToken ::= */
OpenCStyleCommentToken ::= /*
SemiColonToken ::= ;
ColonToken ::= :

```

## ANNEXE 3

```
typedef struct SegmentationControls {
    BOOL    StartSegmentation           : 1 ;
    BOOL    EnablePacingEngine          : 1 ;
    BOOL    EnablePacingClock           : 1 ;
    BOOL    IdleCellGeneration          : 1 ;
} SEGMENTATIONCONTROL ;
```

```

ReadFunction ::= ReturnToken ReadFunctionName ParameterList Function-
Declaration

ReturnToken ::= Ust20wordToken

ReadFunctionName ::= ReadFromToken TypedefStructureName

ParameterList ::= <OpenBracketToken> Int32Type AddressToken<CommaTo-
ken> CharType StarDataToken <CommaToken> Int32Type
NumberOfBytesToken <CommaToken> Int32Type
CycleToken <CommaToken> TypedefIdentifierName
<StarToken> ParameterIdentifierName<CloseBracket-
Token>

FunctionDeclaration ::= <OpenBraceToken> FunctionBody <ClosingBrace-
Token>

FunctionBody ::= PsuedoRegisterDeclaration PsuedoRegisterInitialisa-
tion PsuedoRegisterAssignment TransferWordToTranspu-
terInvocation ReturnStatement

PsuedoRegisterDeclaration ::= Ust20type PsuedoRegisterToken <Semi-
colonToken> <CarriageReturnToken>

PsuedoRegisterInitialisation ::= PsuedoRegisterToken <EqualsToken>
ZeroToken <SemicolonToken> <CarriageRe-
turnToken>

PsuedoRegisterAssignment ::= AttributesToRegisterAssignments

AttributesToRegisterAssignments ::= AttributesToRegisterAssignment
[AttributesToRegisterAssignment]

AttributesToRegisterAssignment ::= PsuedoRegister <BitOrToken>
AttributeDereferenceAndShift <SemicolonTo-
ken> <CarriageReturnToken>

AttributeDereferenceAndShift ::= <OpenBracketToken> <OpenBracketTo-
ken> <OpenBracketToken> <OpenBracketTo-
ken> CastToUst20DataType
<CloseBracketToken> <OpenBracketToken>
AttributeDereference <CloseBracketToken>
<CloseBracketToken> <ShiftUpToken> Attrib-
uteInRegisterBitsShift <CloseBracketTo-
ken> <BitAndToken>
AttributeInRegisterBits <CloseBracketTo-
ken>

CastToUst20DataType ::= <OpenBracketToken> Ust20type <CloseBracketTo-
ken>

AttributeDereference ::= ParameterIdentifierName<ArrowToken>Attrib-
uteIdentifierName

TransferWordToTransputerInvocation ::= TransferWordToTransputerToken
TransferWordToTransputerInvo-
cationParameterList <SemicolonTo-
ken> <CarriageReturnToken>

TransferWordToTransputerInvocationParameterList ::= <OpenBracketTo-
ken> TransferWord-
ToTransputerInvo-
cationParameters
<CloseBracketTo-
ken>

TransferWordToTransputerInvocationParameters ::= DataToken<CommaTo-
ken>AddressOfPsue-
doRegisterToken

ReturnStatement ::= ReturnToken CastToUst20DataType NumberOfBytesTo-
ken <SemicolonToken> <CarriageRe-
turnToken>

```

[illegible][illegible]

## ANNEXE 6

```

WriteFunction ::= WriteReturnType WriteFunctionName WriteParameterList
WriteFunctionName ::= WriteToken Token typedefStructureName
WriteParameterList ::= <OpenBracketToken> Int32Type AddressToken <CommaToken> CharType StarDataToken <CommaToken> Int32Type NumberOfBytesToken <CommaToken> Int32Type CycleToken <CommaToken> TypedefIdentifierName <StarToken> ParameterIdentifierName <CloseBracketToken>
WriteFunctionDeclaration ::= <OpenBraceToken> WriteFunctionBody <ClosingBraceToken>
WriteFunctionBody ::= PsuedoRegisterDeclaration PsuedoRegisterInitialisation TransferWordFromTransputerInvocation RegisterDereferenceAssignment
PsuedoRegisterDeclaration ::= Ust20type PsuedoRegisterToken <SemicolonToken> <CarriageReturnToken>
PsuedoRegisterInitialisation ::= PsuedoRegisterToken <EqualsToken> ZeroToken <SemicolonToken> <CarriageReturnToken>
TransferWordFromTransputerInvocation ::= TransferWordFromTransputerToken TransferWordFromTransputerInvocationParameterList <SemicolonToken> <CarriageReturnToken>
TransferWordFromTransputerInvocationParameterList ::= <OpenBracketToken> TransferWordFromTransputerInvocationParameters <CloseBracketToken>
TransferWordFromTransputerInvocationParameters ::= AddressOfPsuedoRegisterToken <CommaToken> DataToken
RegisterDereferenceAssignment ::= AttributeAssignments
AttributeAssignments ::= AttributeAssignment [AttributeAssignment]
AttributeAssignment ::= AttributeDeferenceAndAssignment
AttributeDeferenceAndAssignment ::= AttributeInRegisterDereference <EqualsToken> ValueAssignment <SemicolonToken> <CarriageReturnToken>
AttributeInRegisterDereference ::= ParameterIdentifierName <ArrowToken> AttributeIdentifierName
ValueAssignment ::= ValueDerivedFromRegister
ValueDerivedFromRegister ::= CastToAttributeType <OpenBracketToken> <OpenBracketToken> PsuedoRegister <BitAndToken> AttributeInRegisterBitsConstant <CloseBracketToken> <ShiftDownToken> AttributeInRegisterBitsShift <CloseBracketToken>
CastToAttributeType ::= <OpenBracketToken> AttributeType <CloseBracketToken>

```

## ANNEXE 7

```
#define SEGMENTATIONCONTROLSTARTSEGMENTATIONBIT 0x1
#define SEGMENTATIONCONTROLENABLEPACINGENGINEBIT 0x2
#define SEGMENTATIONCONTROLENABLEPACINGCLOCKBIT 0x4
#define SEGMENTATIONCONTROLIDLECELLGENERATIONBIT 0x8
#define SEGMENTATIONCONTROLSTARTSEGMENTATIONSHIFT 0x0
#define SEGMENTATIONCONTROLENABLEPACINGENGINESHIFT 0x1
#define SEGMENTATIONCONTROLENABLEPACINGCLOCKSHIFT 0x2
#define SEGMENTATIONCONTROLIDLECELLGENERATIONSHIFT 0x3

void WriteToSegmentationControl (INT32 Address, char *Data, INT32
                                NumberOfBytes, INT32 Cycles,
                                SARSEGMENTATIONCONTROL
                                *SegmentationControlRegister)
{
    UST20WORD    PsuedoRegister ;

    PsuedoRegister = 0 ;
    TransferWordFromTransputer (&PsuedoRegister, Data) ;

    SegmentationControlRegister->StartSegmentation = (BOOL)
        ((PsuedoRegister &
          SEGMENTATIONCONTROLSTARTSEGMENTATIONBIT) >>
          SEGMENTATIONCONTROLSTARTSEGMENTATIONSHIFT) ;

    SegmentationControlRegister->EnablePacingEngine = (BOOL)
        ((PsuedoRegister &
          SEGMENTATIONCONTROLENABLEPACINGENGINEBIT) >>
          SEGMENTATIONCONTROLENABLEPACINGENGINESHIFT) ;

    SegmentationControlRegister->EnablePacingClock = (BOOL)
        ((PsuedoRegister &
          SEGMENTATIONCONTROLENABLEPACINGCLOCKBIT) >>
          SEGMENTATIONCONTROLENABLEPACINGCLOCKSHIFT) ;

    SegmentationControlRegister->IdleCellGeneration = (BOOL)
        ((PsuedoRegister &
          SEGMENTATIONCONTROLCONTROLIDLECELLGENERATIONBIT) >>
          SEGMENTATIONCONTROLCONTROLIDLECELLGENERATIONSHIFT) ;
}
```



ANNEXE 8

```

QueryFunction ::= ReturnToken QueryFunctionName ParameterList Func-
tionDeclaration

ReturnToken ::= AttributeToken

QueryFunctionName ::= AttributeIdentifierName InToken TypedefStruc-
tureName

ParameterList ::= <OpenBracketToken> TypedefIdentifierName <StarTo-
ken> ParameterIdentifierName <CloseBracketToken>

FunctionDeclaration ::= <OpenBraceToken> FunctionBody <ClosingBrace-
Token>

FunctionBody ::= ReturnStatement

ReturnStatement ::= ReturnToken RegisterDereference

RegisterDereference ::= ParameterIdentifierName <DashArrow> AttributeI-
dentifierName <SemicolonToken> <CarriageRe-
turnToken>

TypedefStructureName ::= The name given to the structure of the type-
def

```

## ANNEXE 9

```

UST20WORD ValueInBasepagePtr (BASEPAGEPTR *Register)
{
    return Register->Value ;
}

```

## ANNEXE 10

```

SetFunction ::= SetReturnType SetFunctionName SetParameterList Set-
               FunctionDeclaration

SetReturnType ::= VoidType

SetFunctionName ::= := SetToken AttributeIdentifierName InToken Typede-
               fStructureName

SetParameterList ::= <OpenBracketToken> TypedefIdentifierName <Star-
               Token> ParameterIdentifierName AttributeType Val-
               ueIdentifierName <CloseBracketToken>

SetFunctionDeclaration ::= <OpenBraceToken> SetFunctionBody <Closing-
               BraceToken>

SetFunctionBody ::= RegisterDereferenceAssignment

RegisterDereferenceAssignment ::= RegisterDereference <EqualsToken>
               AttributeIdentifierName <SemicolonTo-
               ken> <CarriageReturnToken>

RegisterDereference ::= ParameterIdentifierName<ArrowToken>Attrib-
               uteIdentifierName <EqualsToken> ValueI-
               dentifierName

```

## ANNEXE 11

```
void SetValueInBasepagePtr (BASEPAGEPTR *Pointer, UST20WORD AnyName)
{
    Pointer->Value = AnyName ;
}
```

## ANNEXE 12

```

ReadFunction ::= ReturnToken ReadFunctionName ParameterList Function-
Declaration

ReturnToken ::= UWordToken

ReadFunctionName ::= ReadFromToken TypedefStructureName

ParameterList ::= <OpenBracketToken> VolatileToken TypedefIdentifi-
erName StarToken ParameterIdentifierName <Close-
BracketToken>

FunctionDeclaration ::= <OpenBraceToken> FunctionBody <ClosingBrace-
Token>

FunctionBody ::= ReturnStatement

ReturnStatement ::= ReturnToken CastToUWordDataType RegisterDerefer-
ence

CastToUWordDataType ::= <OpenBracketToken> UWordtype <CloseBracketTo-
ken>

RegisterDereference ::= ParameterIdentifierName<ArrowToken>Attrib-
uteIdentifierName <SemicolonToken> <Car-
riageReturnToken>

```

## ANNEXE 13

```
#define SEGMENTATIONCONTROLSTARTSEGMENTATIONBIT 0x1

#define SEGMENTATIONCONTROLENABLEPACINGENGINEBIT 0x2

#define SEGMENTATIONCONTROLENABLEPACINGCLOCKIT 0x4

#define SEGMENTATIONCONTROLIDLECELLGENERATIONBIT 0x8

UWORD ReadFromSegmentationControl (volatile SEGMENTATIONCONTROL
                                   *Pointer)

{
    return (UWORD) Pointer->Value ;
}
```

## ANNEXE 15

```
RegisterDereference ::= ParameterIdentifierName<ArrowToken>Attrib-
                        uteIdentifierName <EqualsToken> ValueI-
                        dentifierName <SemicolonToken>
                        <CarriageReturnToken>
```

```
{
    Pointer->Value = Value ;
}
```

```

Documentation ::= FontCatalog TableDefinitions Paragraphs

FontCatalog ::= OpenStatement Fonts CloseStatement
Fonts ::= Font1 Font2 Font3
Font1 ::= OpenStatement FontToken Tag1 Family Angle Weight1 Size1 CloseStatement
Font2 ::= OpenStatement FontToken Tag2 Family Angle Weight2 Size1 CloseStatement
Font3 ::= OpenStatement FontToken Tag3 Family Angle Weight1 Size2 CloseStatement
Family ::= OpenStatement FamilyToken OpenSingleQuote FamilyString CloseSingleQuote CloseStatement
Angle ::= OpenStatement AngleToken OpenSingleQuote AngleString CloseSingleQuote CloseStatement
Tag1 ::= OpenStatement TagToken OpenSingleQuote Tag1String CloseSingleQuote CloseStatement
Tag2 ::= OpenStatement TagToken OpenSingleQuote Tag2String CloseSingleQuote CloseStatement
Tag3 ::= OpenStatement TagToken OpenSingleQuote Tag3String CloseSingleQuote CloseStatement
Weight1 ::= OpenStatement WeightToken OpenSingleQuote Weight1String CloseSingleQuote CloseStatement
Weight2 ::= OpenStatement WeightToken OpenSingleQuote Weight2String CloseSingleQuote CloseStatement
Size1 ::= OpenStatement SizeToken OpenSingleQuote Size1String CloseSingleQuote CloseStatement
Size2 ::= OpenStatement SizeToken OpenSingleQuote Size2String CloseSingleQuote CloseStatement

TableDefinitions ::= OpenStatement TblsToken TableDefinition {TableDefinition} CloseStatement
TableDefinition ::= OpenStatement TblToken TableConfig TableTitle TableHeader TableBody CloseStatement
TableConfig ::= TableID TableFormat TableNumberColumn TableColumnWidth
TableID ::= OpenStatement TblIDToken IDnumber CloseStatement
TableFormat ::= OpenStatement TblTagToken OpenSingleQuote FormatString CloseSingleQuote CloseStatement
TableNumberColumn ::= OpenStatement TblNumColumnsToken ColumnNumber CloseStatement
TableColumnWidth ::= Width1 Width1 Width2 Width2 Width1 Width3
Width1 ::= OpenStatement TblNumColumnWidthToken Width1String CloseStatement
Width2 ::= OpenStatement TblNumColumnWidthToken Width2String CloseStatement
Width3 ::= OpenStatement TblNumColumnWidthToken Width3String CloseStatement

TableTitle ::= OpenStatement TblTitleToken TableContent CloseStatement
TableContent ::= OpenStatement TblTitleContentToken ParaTitle CloseStatement
ParaTitle ::= OpenStatement ParaToken ParaTitleTag ParaTitleFont ParaTitleLine CloseStatement
ParaTitleTag ::= OpenStatement PgffTagToken OpenSingleQuote PgffTagNameString CloseSingleQuote CloseStatement

```





```

Cell6ContentBody ::= OpenStatement CellContentToken Para6Body CloseS-
                    tatement
Para1Body ::= OpenStatement ParaToken ParaBodyTag ParaBodyFont
              ParaLine1Body CloseStatement
ParaLine1Body ::= OpenStatement ParaLineToken String1Body CloseState-
                  ment
String1Body ::= OpenStatement StringToken OpenSingleQuote Body1Name
                CloseSingleQuote CloseStatement
Para2Body ::= OpenStatement ParaToken ParaBodyTag ParaBodyFont
              ParaLine2Body CloseStatement
ParaLine2Body ::= OpenStatement ParaLineToken String2Body CloseState-
                  ment
String2Body ::= OpenStatement StringToken OpenSingleQuote Body2Name
                CloseSingleQuote CloseStatement
Para3Body ::= OpenStatement ParaToken ParaBodyTag ParaBodyFont
              ParaLine3Body CloseStatement
ParaLine3Body ::= OpenStatement ParaLineToken String3Body CloseState-
                  ment
String3Body ::= OpenStatement StringToken OpenSingleQuote Body3Name
                CloseSingleQuote CloseStatement
Para4Body ::= OpenStatement ParaToken ParaBodyTag ParaBodyFont
              ParaLine4Body CloseStatement
ParaLine4Body ::= OpenStatement ParaLineToken String4Body CloseState-
                  ment
String4Body ::= OpenStatement StringToken OpenSingleQuote Body4Name
                CloseSingleQuote CloseStatement
Para5Body ::= OpenStatement ParaToken ParaBodyTag ParaBodyFont
              ParaLine5Body CloseStatement
ParaLine5Body ::= OpenStatement ParaLineToken String5Body CloseState-
                  ment
String5Body ::= OpenStatement StringToken OpenSingleQuote Body5Name
                CloseSingleQuote CloseStatement
Para6Body ::= OpenStatement ParaToken ParaBodyTag ParaBodyFont
              ParaLine6Body CloseStatement
ParaLine6Body ::= OpenStatement ParaLineToken String6Body CloseState-
                  ment
String6Body ::= OpenStatement StringToken OpenSingleQuote Body6Name
                CloseSingleQuote CloseStatement
ParaBodyFont ::= OpenStatement PgffontToken TagBody CloseStatement

TagBody ::= OpenStatement TagToken OpenSingleQuote TagBodyString Clos-
            eSingleQuote CloseStatement
ParaBodyTag ::= OpenStatement PgffontToken OpenSingleQuote PgffontTagBodyS-
                tring CloseSingleQuote CloseStatement

Paragraphs ::= Paragraph {Paragraph}
Paragraph ::= OpenStatement ParaToken ParagraphContent CloseStatement
ParagraphContent ::= ParagraphTag ParagraphLine
ParagraphTag ::= OpenStatement PgffontToken OpenSingleQuote BODY Clos-
                eSingleQuote CloseStatement
ParagraphLine ::= OpenStatement ParaLineToken ParaLineContent Clos-
                  tatement
ParaLineContent ::= OpenStatement ATblToken IDNumber CloseStatement

```

```

TagHeaderString ::= Tag1String
TagTitleString ::= Tag3String
TagBodyString ::= Tag2String
PgftagTitleString ::= TableTitle
PgftagHeaderString ::= CellHeading
PgftagBodyString ::= CellBody
Width1String ::= 0.7
Width2String ::= 1.5
Width3String ::= 1.0
FormatString ::= Format A
FamilyString ::= Times
AngleString ::= Regular
Tag1String ::= NoWareHeading
Tag2String ::= NoWareBody
Tag3String ::= NowareTitle
Weight1String ::= Bold
Weight2String ::= Regular
Size1String ::= 11.0 pt
Size2String ::= 12.0 pt
OpenSingleQuote ::= `
CloseSingleQuote ::= `

```



```

TableName ::= TypedefIdentifierName
Header1Name ::= Word Offset
Header2Name ::= Bit Offset
Header3Name ::= Bit Field
Header4Name ::= Function
Header5Name ::= Reset State
Header6Name ::= R/W
Body1Name ::= Current register number
Body2Name ::= Size of the attribute given in the input data structure
Body3Name ::= Name of the attribute given in the input data structure
Body4Name ::= Comment associated with the attribute within the input
data structure
Body5Name ::= 0 | NULL
Body6Name ::= R/W comment given in the input data strucure

IDnumber ::= start at 1 and increment it
ColumnNumber ::= 6
StraddleNumber ::= Number of attributes within the current register
TableIdentifierName ::= Name given to the typedef of the data structure

CloseStatement ::= >
OpenStatement ::= <

```

What is claimed is:

1. A method of operating a computer system to design an application specific processor (ASP) comprising:

defining a set of peripherals for the ASP which are responsive to stimuli and which communicate with a processor;

generating for each peripheral an input file which defines the functional attributes of that peripheral in a high level language with an input data structure;

entering the input file into the computer system and operating a modelling tool loaded on the computer system to generate from the input file a register definition file by allocating specific elements of the input data structure to predefined sectors of a register definition table; and

using the register definition file to create in silicon the registers of the ASP.

2. A method according to claim 1, wherein each input file comprises a data structure which defines for each of a set of registers the name of an element in the register, the bit length of the element, the functional status of the element and the function of the element.

3. A method according to claim 1, wherein each register definition table includes at least predefined sectors for the bit location within a register of an element, the name of the element, the function of the element and the functional status of the element.

4. A method according to claim 1, wherein the register definition table includes the word location of the register within a memory map for access during simulation of the ASP.

5. A computer system which comprises a processor and a memory, the memory holding a program representing a modelling tool for use in designing an application specific processor (ASP), wherein the computer system comprises an input means for receiving a

the processor being operable to execute the program representing the modelling tool to generate from the input file a register definition file by allocating specific elements of the input data structure to predefine sectors of a register definition table; and

6. A computer system according to claim 5, wherein the input means comprises means for receiving a physical recording device holding the input file for each peripheral.

8. A computer program product stored on a computer readable medium and comprising software code portions operable when executed by a computer to read an input file which defines in an input data structure the functional attributes of a peripheral for an application specific processor in a high level language, and to generate from that input file a register definition file, the software code portions including a code portion for allocating specific elements of the input data structure to predefined sectors of a register definition table for each of a plurality of registers.

9. A register definition file stored on a computer readable medium and comprising a plurality of register definition tables, each table including at least predefined sectors for the bit location within a register of an element, the name of the element, the function of the element and the functional status of the element, and each table further including the word

location of the register within a memory map for access during simulation of an ASP implementing the registers.

668250 4444450

Abstract of the DisclosureDESIGN OF AN APPLICATION SPECIFIC PROCESSOR (ASP)

A method of operating a computer system to design an application specific processor (ASP) comprises defining a set of peripherals for the ASP which are responsive to stimuli and which communicate with a processor, generating for each peripheral an input file which defines the functional attributes of that peripheral in a high level language with an input data structure, entering the input file into the computer system and operating a modelling tool loaded on the computer system to generate from the input file a register definition file by allocating specific elements of the input data structure to predefined sectors of a register definition table, and using the register definition file to create in silicon the registers of the ASP.

a b c d e f g h i j k l m n o p q r s t u v w x y z

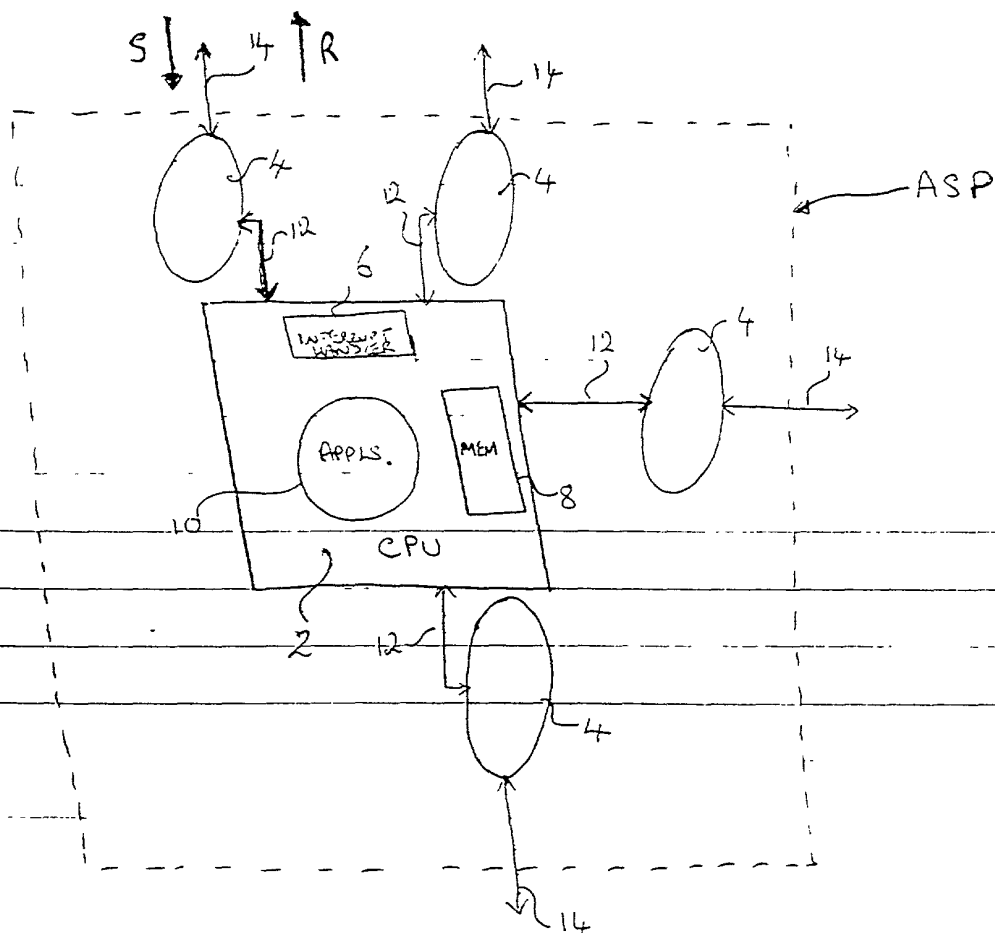


Figure 1

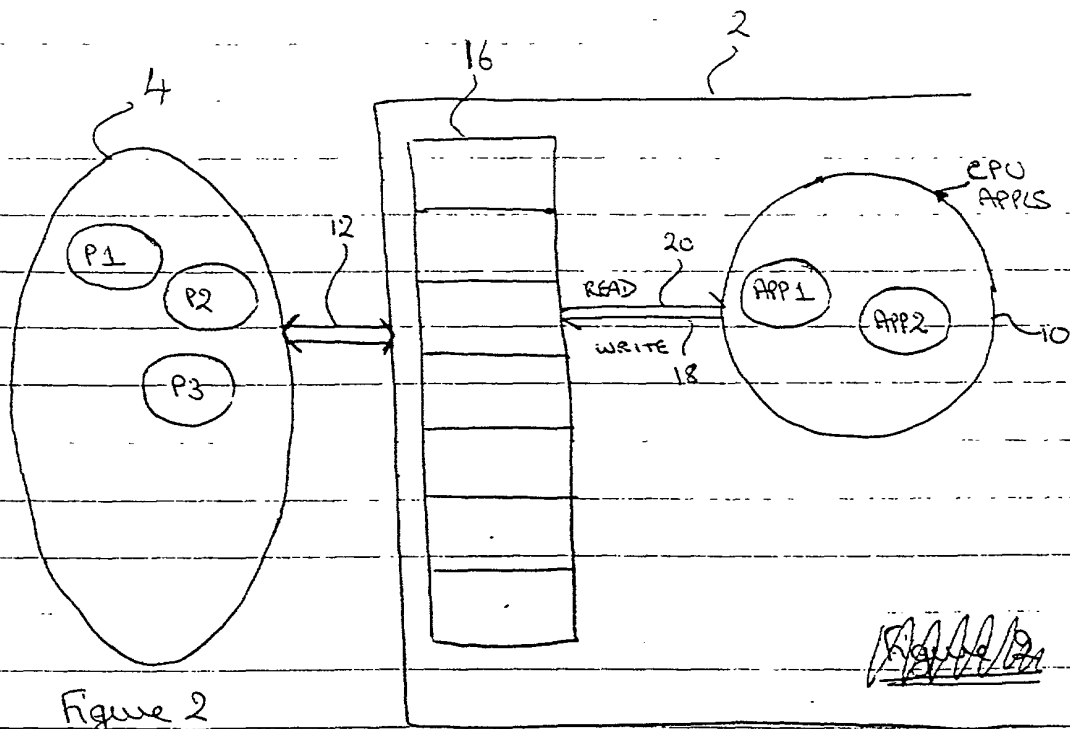


Figure 2

*[Handwritten signature]*

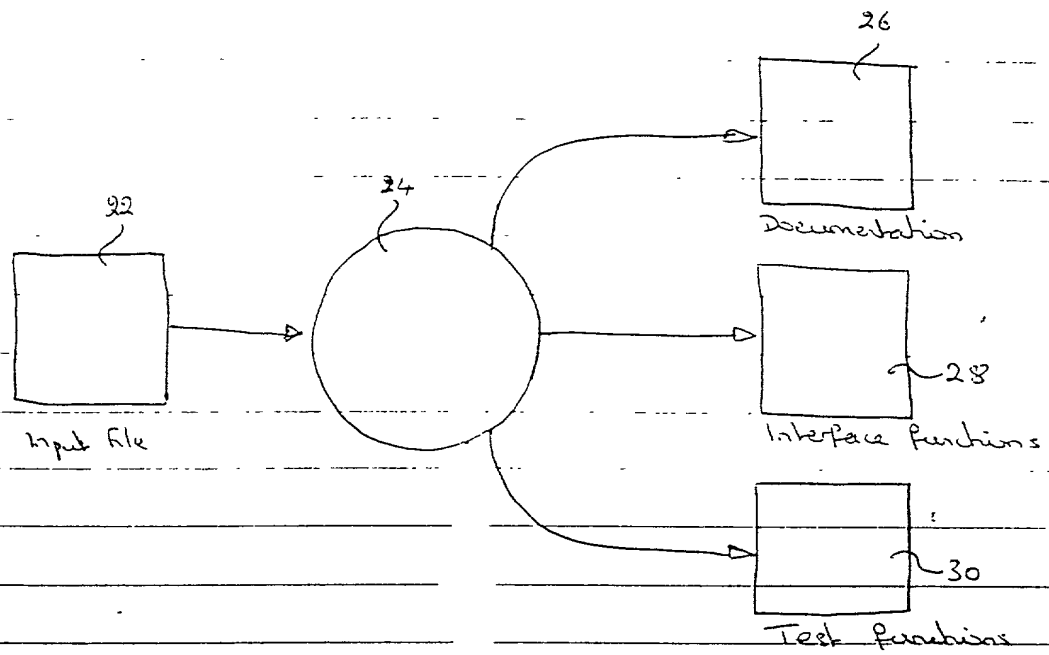
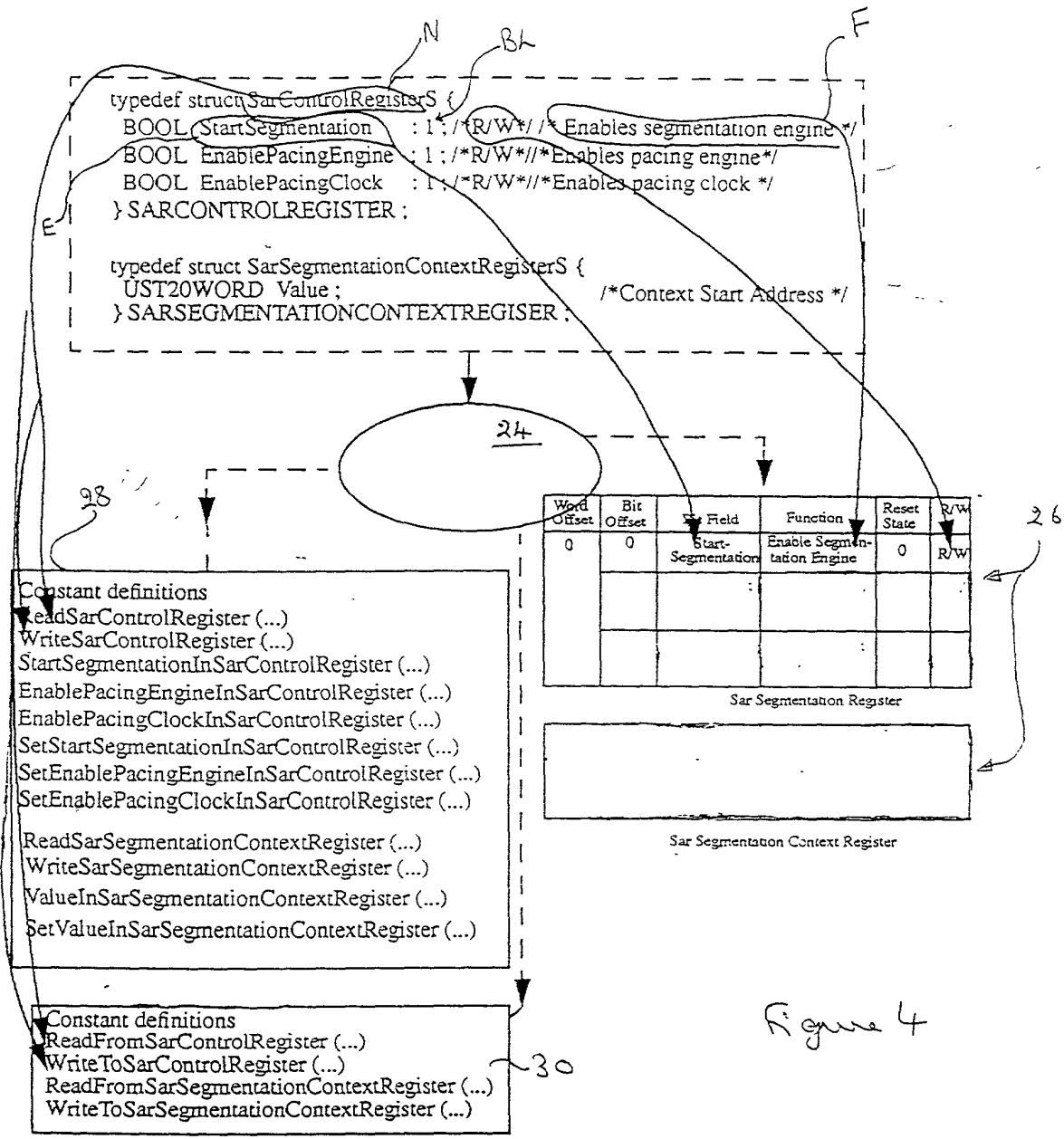
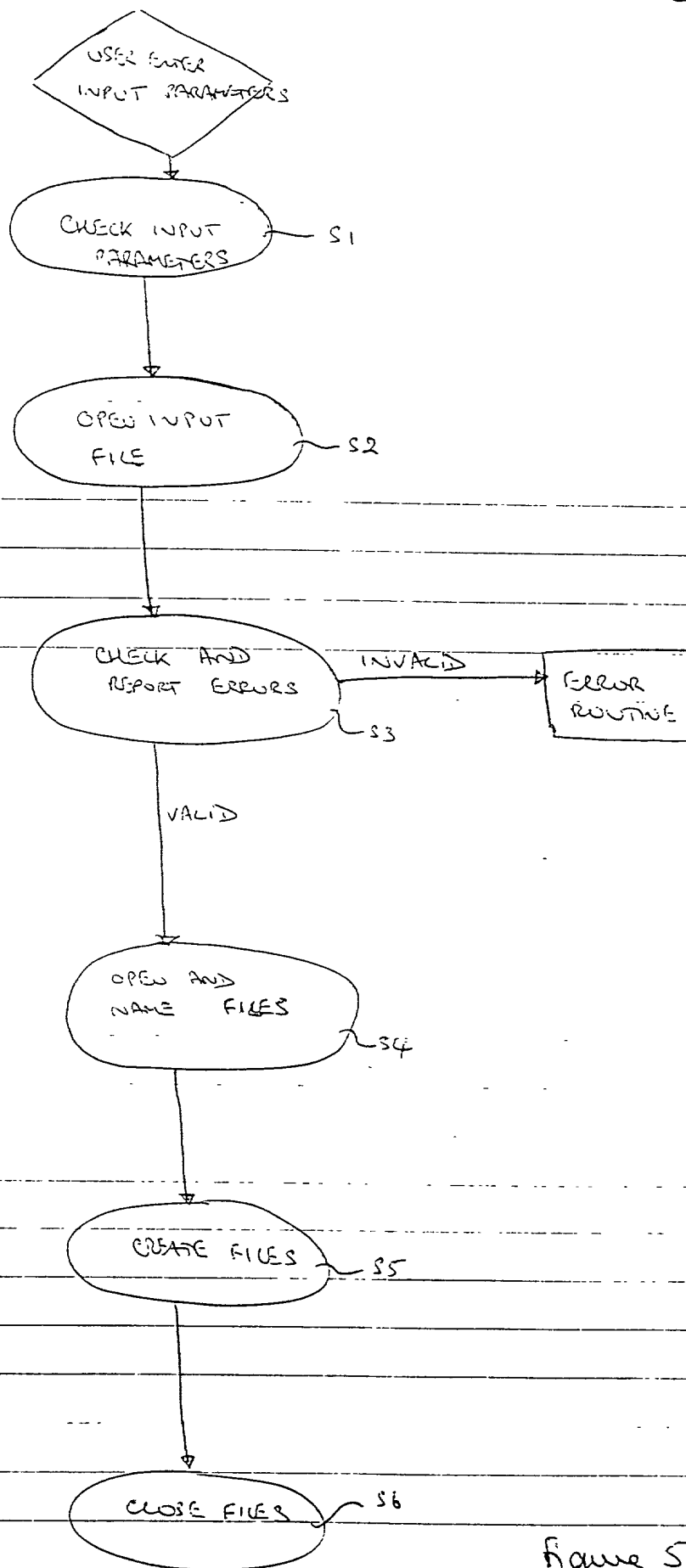


Figure 3





Figure 5

